

# Debugging of Memory Safety Errors in Memory Managers

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# The Problem

- I set my GC header bit to zero.
- Other runtime system code runs.
- My GC header bit is now set to one.

## The Problem: Errors

- I set my GC header bit to zero.
- Other runtime system code runs.
- My GC header bit is now set to one.
- Memory safety errors
  - Dynamic semantics

## Not The Problem: Bugs

- Statically possible **bugs**
- Infeasible without intervention
- Memory safety bugs:
  - Static semantics

# Printf Debugging and Code Editors

```
72 public class JavaHeader implements JavaHeaderConstants {
73
74     protected static final int SCALAR_HEADER_SIZE = JAVA_HEADER_BYTES + OTHER_HEADER_BYTES;
75     protected static final int ARRAY_HEADER_SIZE = SCALAR_HEADER_SIZE + ARRAY_LENGTH_BYTES;
76
77     /** offset of object reference from the lowest memory word */
78     public static final int OBJECT_REF_OFFSET = ARRAY_HEADER_SIZE; // from start to ref
79     protected static final Offset TIB_OFFSET = JAVA_HEADER_OFFSET;
80     protected static final Offset STATUS_OFFSET = TIB_OFFSET.plus(STATUS_BYTES);
81     public static final Offset AVAILABLE_BITS_OFFSET =
82         VM.LittleEndian ? (STATUS_OFFSET) : (STATUS_OFFSET.plus(STATUS_BYTES - 1));
83 }
```

- Mental math

- Source code modification

- Static control flow graph

- Mailing lists

# Documentation to the rescue!

```
61 * |<- lo memory                                     hi memory ->|
62 *
63 *   SCALAR LAYOUT:
64 * |<----- scalar header ----->|
65 * +-----+-----+-----+-----+-----+-----+
66 * | GCHheader | MiscHeader | JavaHeader | fld0 | fld1 | fldx | fldN-1 |
67 * +-----+-----+-----+-----+-----+-----+
68 *                ^ JH0FF                ^objref
69 *
```

# The Current Solution

- Weeks or more of manual debugging efforts
- False assumption: all program errors are the result of **bugs**
  - Unspecified behavior need not be a program bug
  - Documentation is a PL concern
- A program error is a developer ascribing correctness to behavior

# PL Approach: What information do we need?

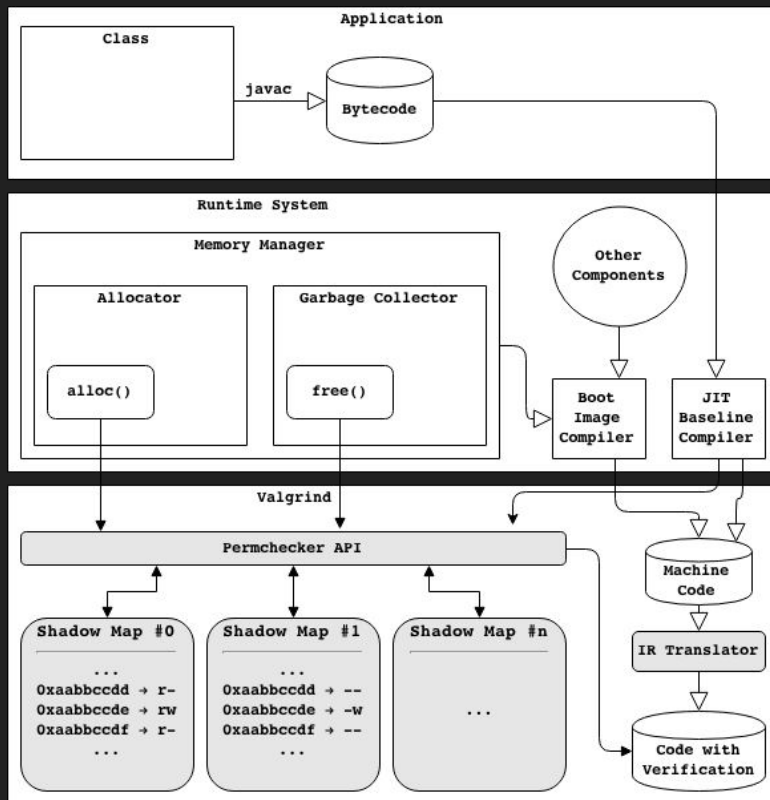
- What's the current configuration of memory?
- What code is running?
- What memory is that code allowed to access?

# PL Approach: Every problem is a language

- What's the current configuration of memory?
- What code is running?
- What memory is that code allowed to access?
- $M(\text{VM}) = \{ \text{MemoryConfig} \}$ 
  - Is  $(0x0000 \dots 0xFFFF) \in M(\text{VM})?$
- $C(\text{VM}) = \{ (\text{Code}, \text{MemoryType}) \}$ 
  - Is  $(\text{Code}(ip), \text{MemoryType}(esi)) \in C(\text{VM})?$

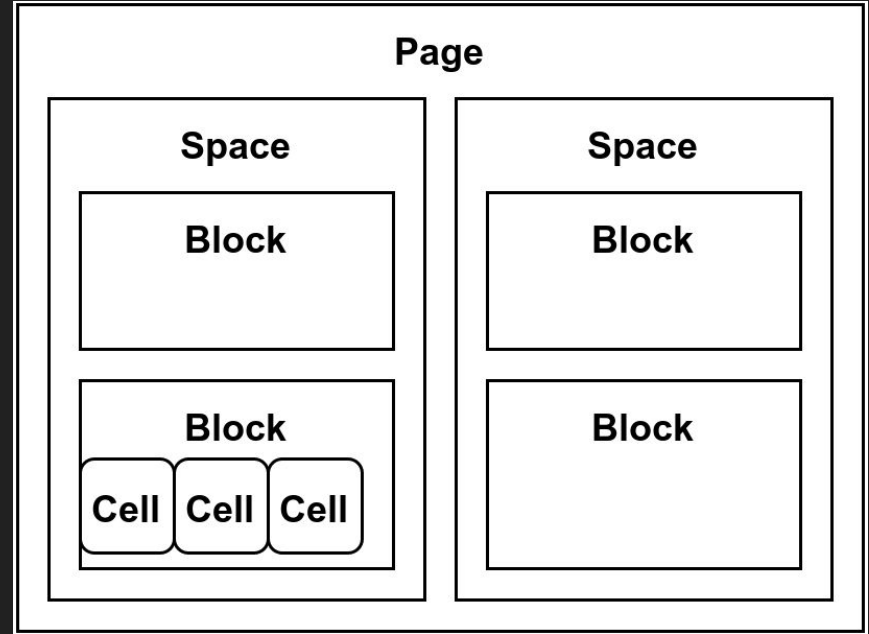
# Prior Work: Dynamic Binary Instrumentation

- Specify M(VM) and C(VM) imperatively "at the same time"
- Half portable
  - VM-specific alloc() and free()
  - Shared Permchecker API



# Ongoing Work: Specify M(VM) Separately

- MM memory is layered
  - Page → Space → Block → Cell
- MM memory is bifurcated
  - (1) alloc, and (2) free
- MM memory is "owned"
  - Metadata at each layer
  - Free chunks at each layer



# Ongoing Work: Specification Language (FLP)

- MM memory is layered
  - Page → Space → Block → Cell
- MM memory is bifurcated
  - (1) alloc, and (2) free
- MM memory is "owned"
  - Metadata at each layer
  - Free chunks at each layer

```
Heap |heapSz pages| @(1 page) heapSz ->
( mapped :
  ( # Region |2^22 bytes| @(2^22 bytes) ->
    ( Meta |sz pages| ->
      ( next : Ptr Block
        , bmd : 1 byte
        , csc : 1 byte
        , iu   : 2 bytes
        , fl_meta : Ptr Block
        , pad1 bytes)
      , # Block |sz pages| @(1 page) ->
        ( Free ->
          ( prevBlock : Ptr Free
            // More fields go here
            , pad2 bytes)
          | Alloc ->
            ( Bump ->
              ( objs : # Object
                , free : # bytes)
              | LargeObj -> # bytes))))
        , unmapped : # bytes)
```

# Ongoing Work: Shared Language (FLP)

```
Heap m k ->
# MegaBlock |2^m bytes| @(2^m bytes) ->
// -----Megablock "Header"-----
( Descrs |2^k bytes| ->
  ( padMB bytes
    // ++++++
  , bds : n BlockDescr |2^d bytes| @(2^k bytes) ->
    ( start : Ptr Stg.Word
      , free : Ptr Stg.Word
      , link : Ptr BlockDescr
      , ( back : Ptr BlockDescr
        | bitmap : Ptr Stg.Word
        | scan : Ptr Stg.Word)
      , gen : Ptr Generation
      , gen_no : Stg.Word16
      , dest_no : Stg.Word16
      , node : Stg.Word16
      , Flags |Stg.Word16| ->
        ( LARGE | EVACUATED | FREE
          | PINNED | MARKED | KNOWN
          | EXEC | FRAGMENTED | SWEPT
          | COMPACT )
      , n_blocks : Stg.Word32
      , padD bytes)))
// ++++++
// -----Megablock payload-----
, blocks : n Block |2^k bytes| @(2^k bytes) ->
  ( closures : # Stg.Closures
    , free : # bytes)
```

```
Heap |heapSz pages| @(1 page) heapSz ->
( mapped :
  ( # Region |2^22 bytes| @(2^22 bytes) ->
    ( Meta |sz pages| ->
      ( next : Ptr Block
        , bmd : 1 byte
        , csc : 1 byte
        , iu : 2 bytes
        , fl_meta : Ptr Block
        , pad1 bytes)
      , # Block |sz pages| @(1 page) ->
        ( Free ->
          ( prevBlock : Ptr Free
            // More fields go here
            , pad2 bytes)
          | Alloc ->
            ( Bump ->
              ( objs : # Object
                , free : # bytes)
              | LargeObj -> # bytes))))
      , unmapped : # bytes)
```

# Ongoing Work: Satisfying constraints

```
Heap m k ->
# MegaBlock |2^m bytes| @(2^m bytes) ->
// -----Megablock "Header"-----
( Descrs |2^k bytes| ->
  ( padMB bytes
    // ++++++
  , bds : n BlockDescr |2^d bytes| @(2^k bytes) ->
    ( start : Ptr Stg.Word
      , free : Ptr Stg.Word
      , link : Ptr BlockDescr
      , ( back : Ptr BlockDescr
        | bitmap : Ptr Stg.Word
        | scan : Ptr Stg.Word)
      , gen : Ptr Generation
      , gen_no : Stg.Word16
      , dest_no : Stg.Word16
      , node : Stg.Word16
      , Flags |Stg.Word16| ->
        ( LARGE | EVACUATED | FREE
          | PINNED | MARKED | KNOWN
          | EXEC | FRAGMENTED | SWEPT
          | COMPACT )
      , n_blocks : Stg.Word32
      , padD bytes)))
// ++++++
// -----Megablock payload-----
, blocks : n Block |2^k bytes| @(2^k bytes) ->
  ( closures : # Stg.Closures
    , free : # bytes)
```

$$\begin{aligned} 2^d &= \text{sz}(\text{BlockDescr}) \\ &= 5 \times \text{sz}(\text{Ptr}) + 4 \times \text{sz}(\text{Word16}) + \\ &= \text{sz}(\text{Word32}) + \text{padD bytes} \\ &= 5 \text{ words} + (12 + \text{padD}) \text{ bytes} \\ d_{32} &= \log_2(32 + \text{padD}) = 5 \text{ [ padD = 0 bytes]} \\ d_{64} &= \log_2(52 + \text{padD}) = 6 \text{ [ padD = 12 bytes]} \end{aligned}$$

- General constraint solving

# Ongoing Work: FLP Toolchain

- Parser implemented in Haskell
- Currently implementing transpiler
  - Tile FLP "layers" onto C structs or C++ classes
  - Examples: Hotspot, Jikes RVM, GHC
- Pin and Valgrind DBI backend
- Take-away: a language for specifying memory layout capable of describing a wide variety of VM implementations (for Java, Haskell, Javascript, ...)

# Future Work: Specifying C(VM)

- Integrate with host language or a framework, options include:
  - Target C++ class methods (portability: lowest common denominator)
  - Target Rust (safety: stronger guarantees)
  - Integrate with LLVM (automation)
  - Other suggestions?